

Motor: Enabling Multi-Versioning for Distributed Transactions on

Disaggregated Memory

Ming Zhang, Yu Hua, Zhijun Yang Huazhong University of Science and Technology, China

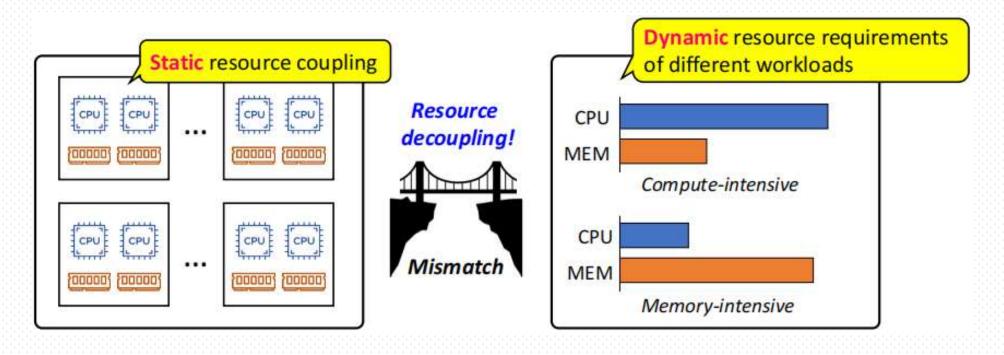
Presented by Sen Han



Background-DM



Insufficient Memory Utilization in Cloud (20%~60%^[1-4])

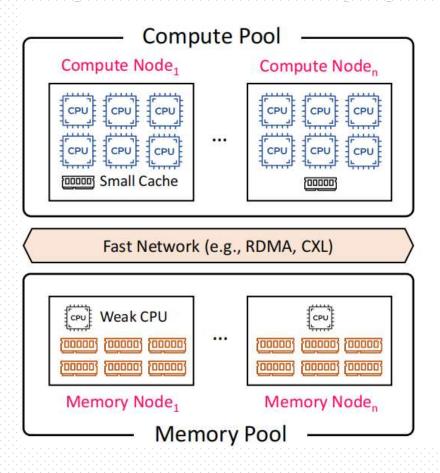


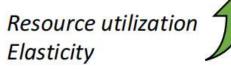
[1] MemTrade@SIGMERTICS'23, Borg@Eurosys'20, LegoOS@OSDI'18
[2] Google Production Cluster Trace. https://github.com/google/cluster-data
[3] Alibaba Production Cluster Trace. https://github.com/alibaba/clusterdata
[4] Snowflake Dataset. https://github.com/resource-disaggregation/snowset

Background-DM



Disaggregated Memory(DM)

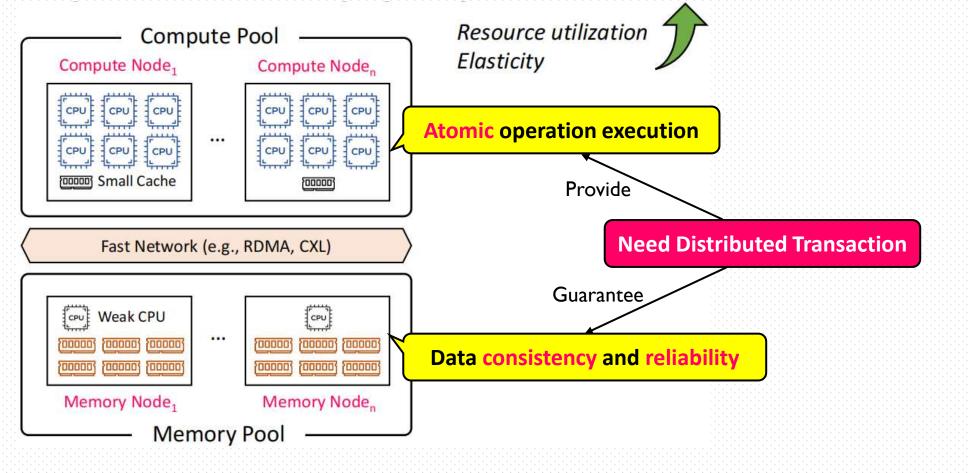




Background-DM

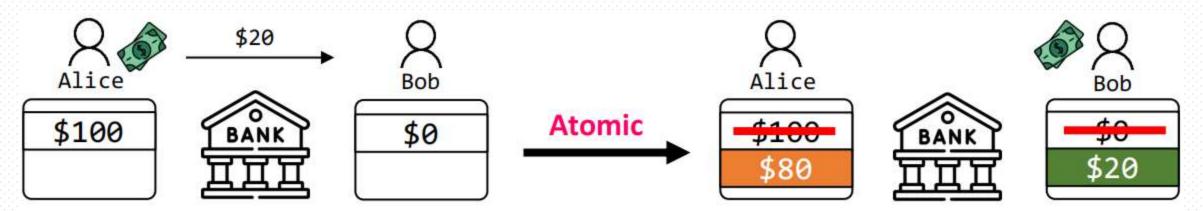


Disaggregated Memory(DM)



Background-Transaction

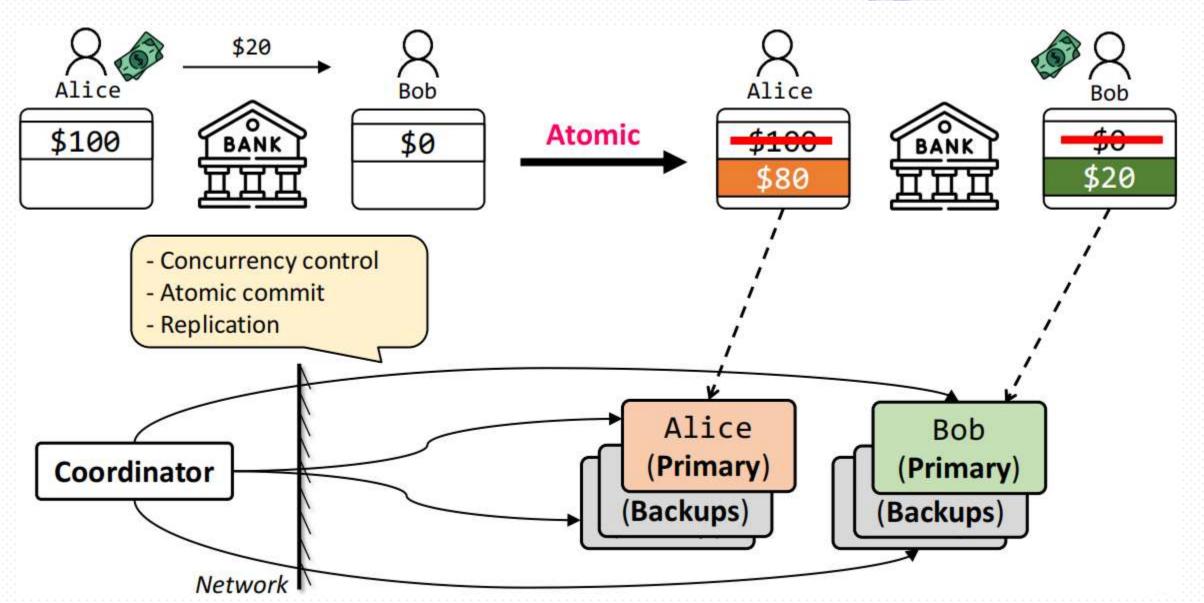
ADSLAB



Txn begi	n		
Alice:	\$100	\rightarrow	\$80
Bob:	\$0	\rightarrow	\$20
Txn end			

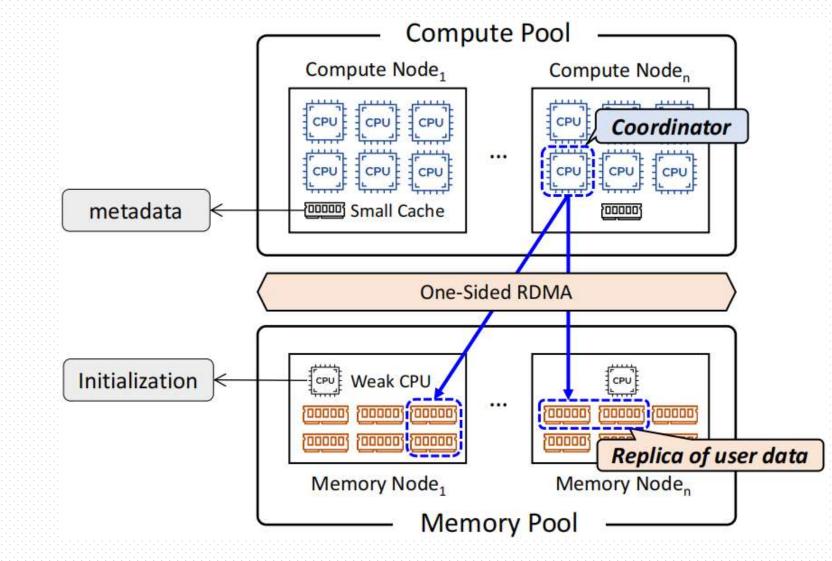
Transaction

Background-Transaction



Background-Transaction on DM

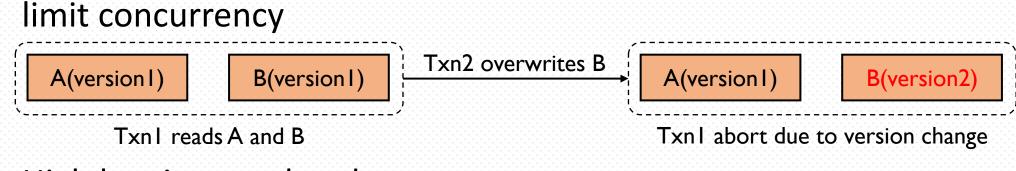




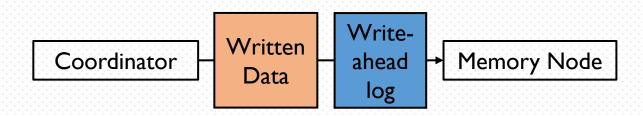




Single-versioning distributed transaction system for DM^[1]



High logging overhead

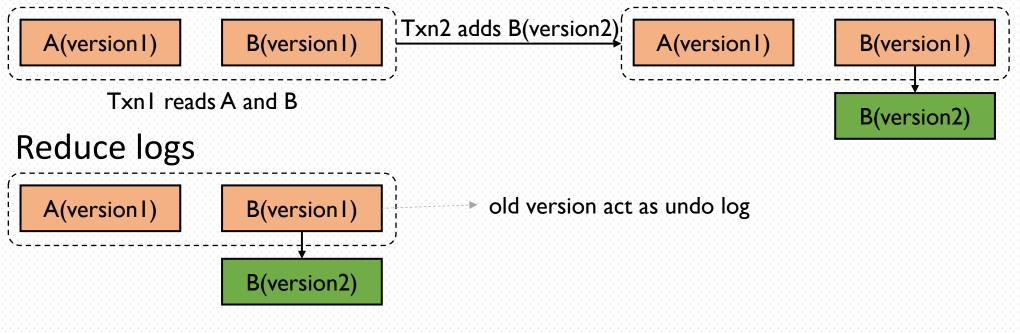






Multi-versioning helps address limitations of single-versioning

Allow more concurrency

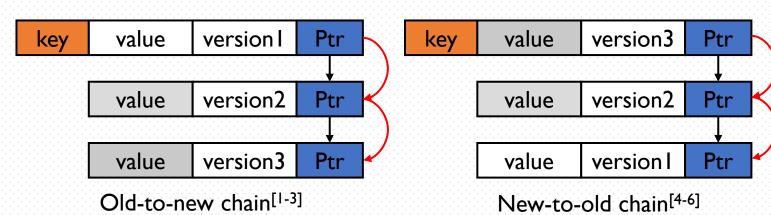


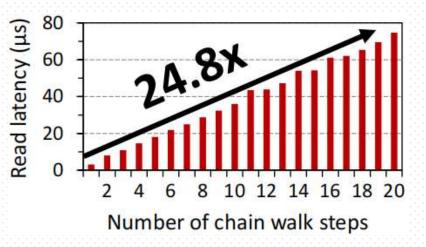
[1] FORD@FAST'22

Does Multi-Versioning Work on DM?

Existing systems are based on monolithic servers

Inefficient linked version chain does not fit DM

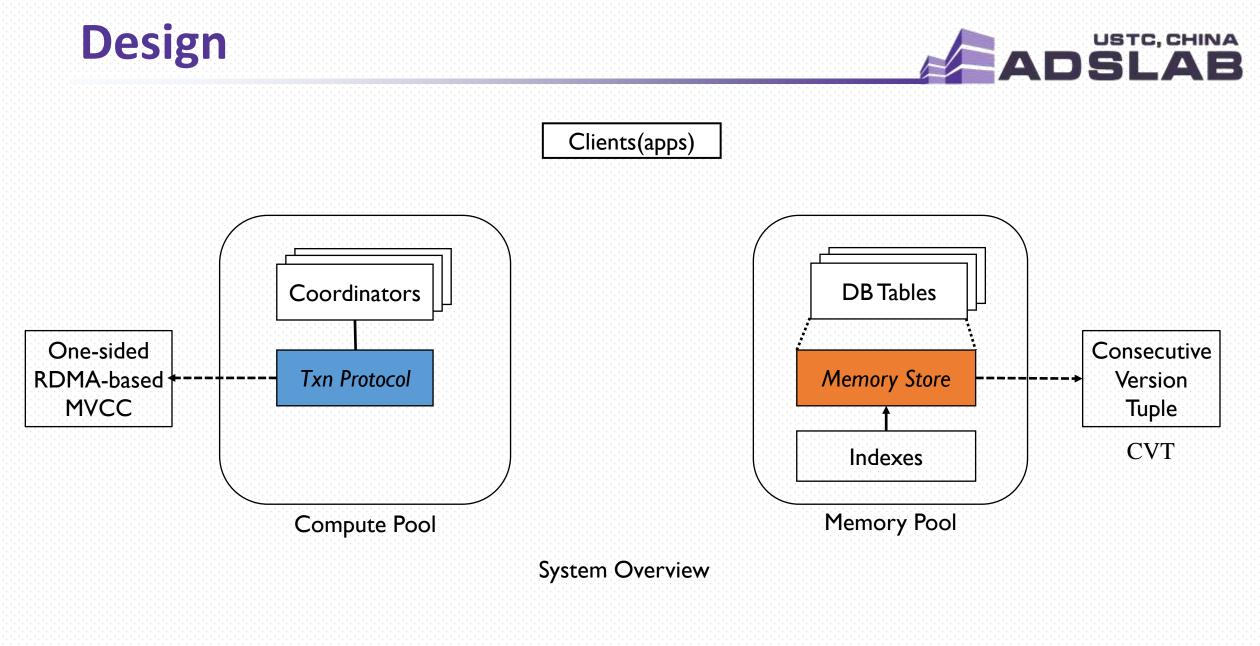


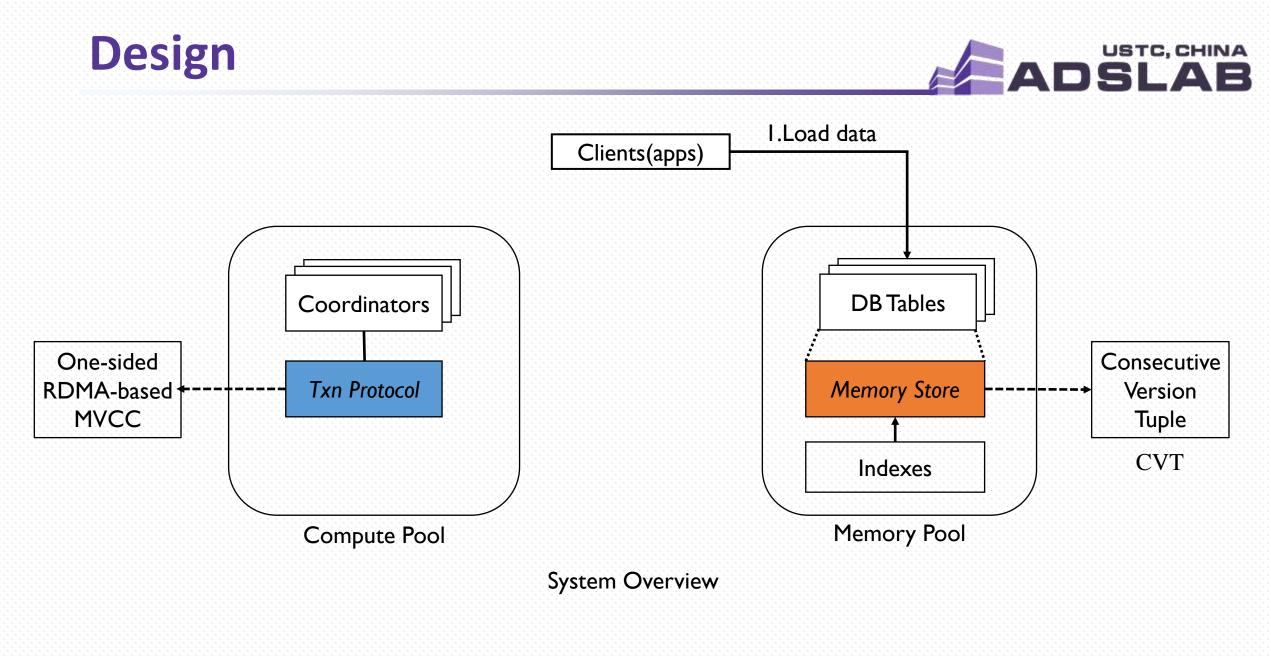


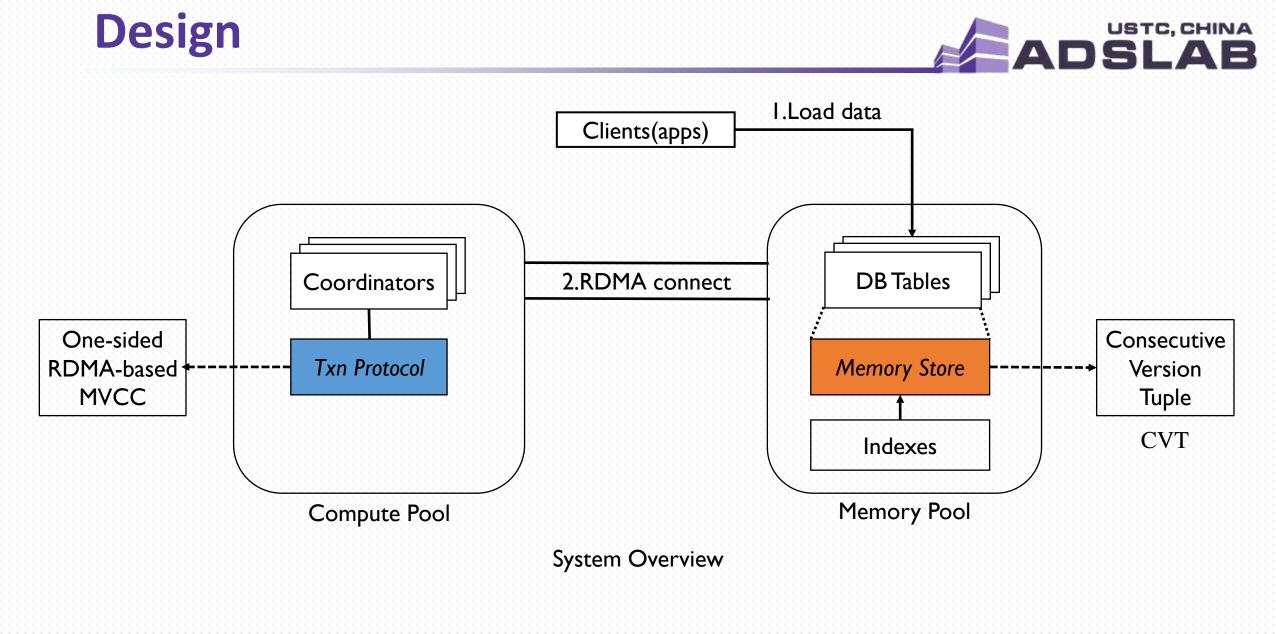
Incompatible transaction protocol

- Frequently consumes CPU of each data node^[1,4,5]
- Memory node stores data but only has weak CPU

[1] DST@NSDI'21 [2] Hekaton@SIGMOD'13 [3] Aurogon@FAST'22[4] FaRMv2@SIGMOD'19 [5] Neumann et al.@SIGMOD'15 [6] NAM-DB@VLDB'17

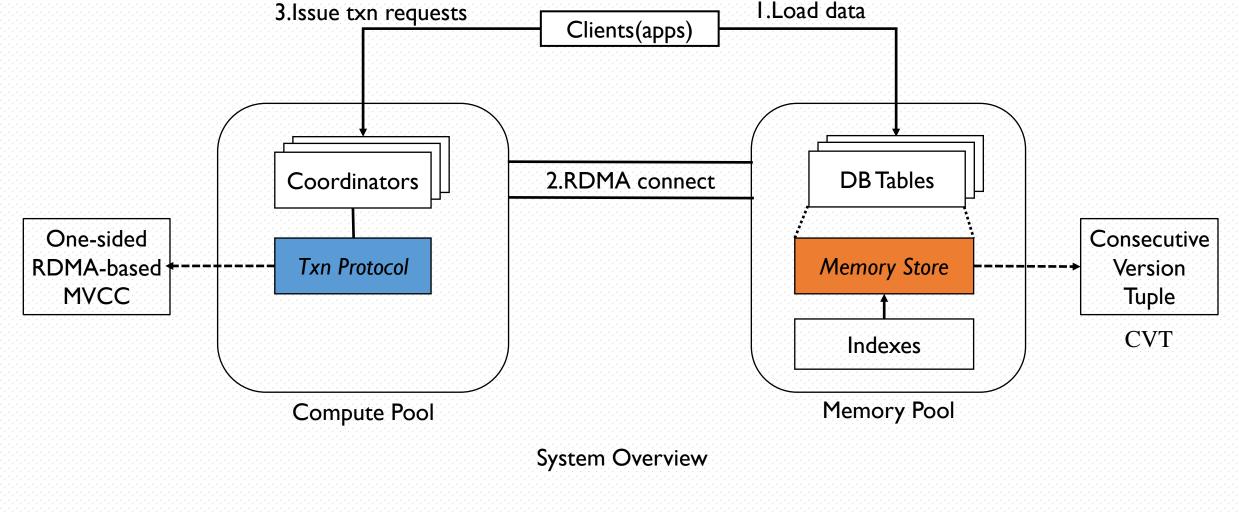






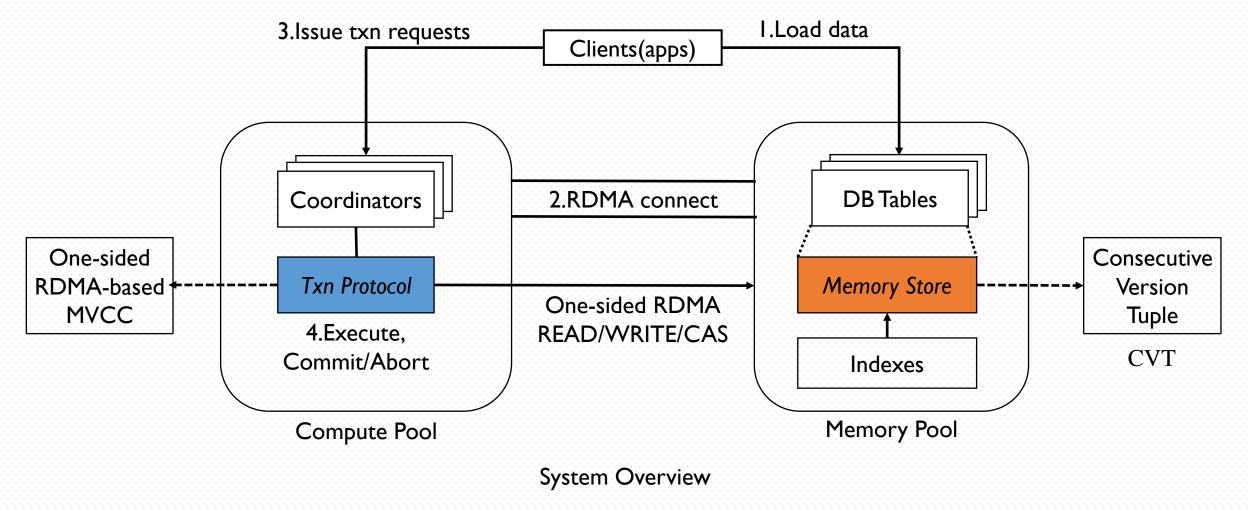






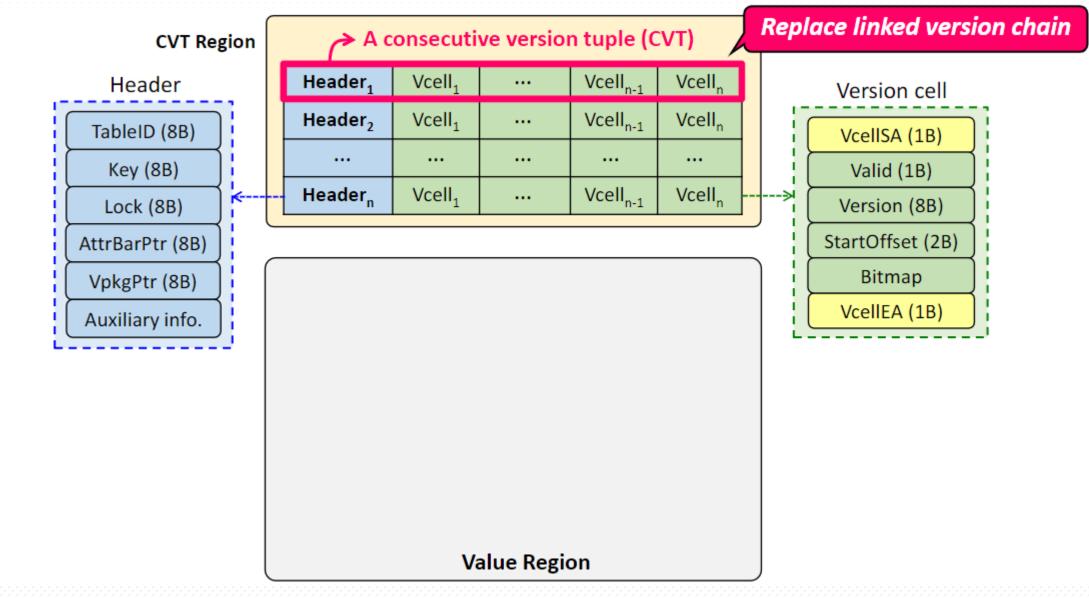






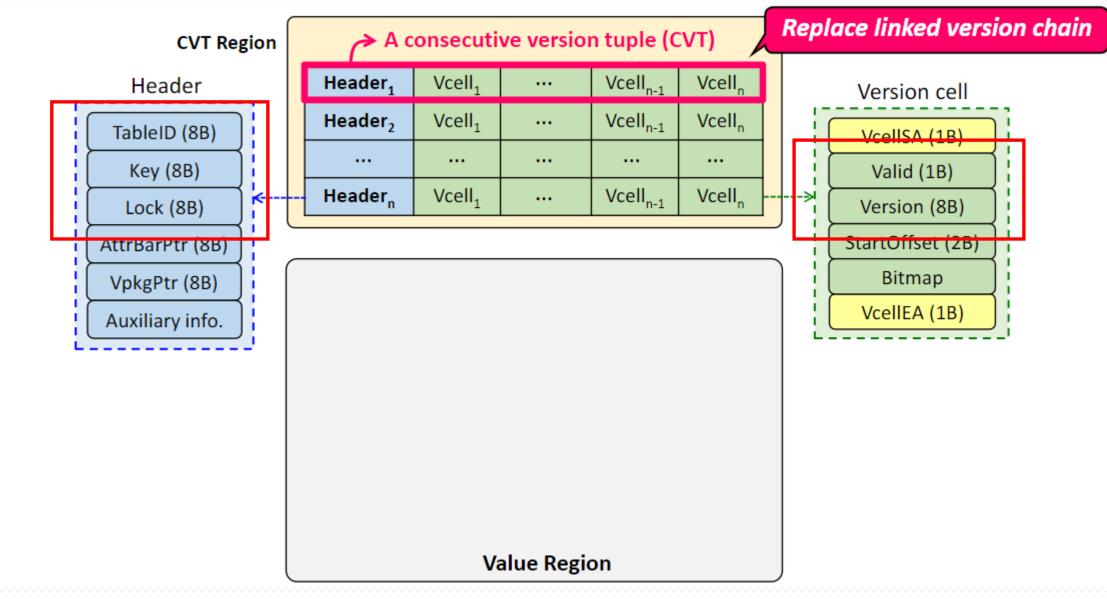
Consecutive Version Tuple





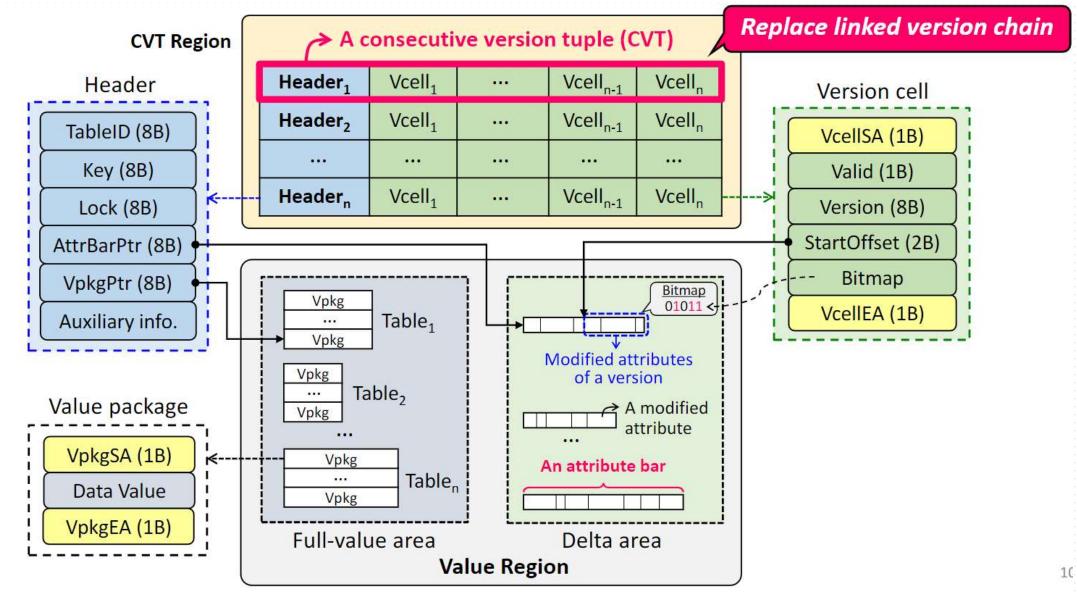
Consecutive Version Tuple





Reducing Memory Overhead

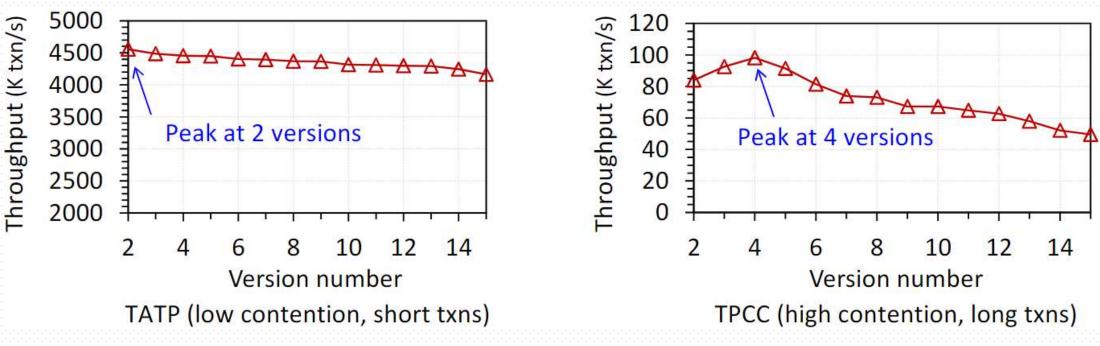




Number of versions to store

Number of Versions: depending on workload characteristics

- Read-write contention
- Number of accessed Records

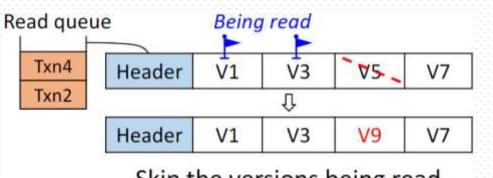


Coordinator-Active Garbage Collection

A CVT runs out of space - GC required

Prior systems track transaction states

• CPU in memory nodes is too weak to frequently track



Skip the versions being read High overhead for compute nodes to maintain remote states

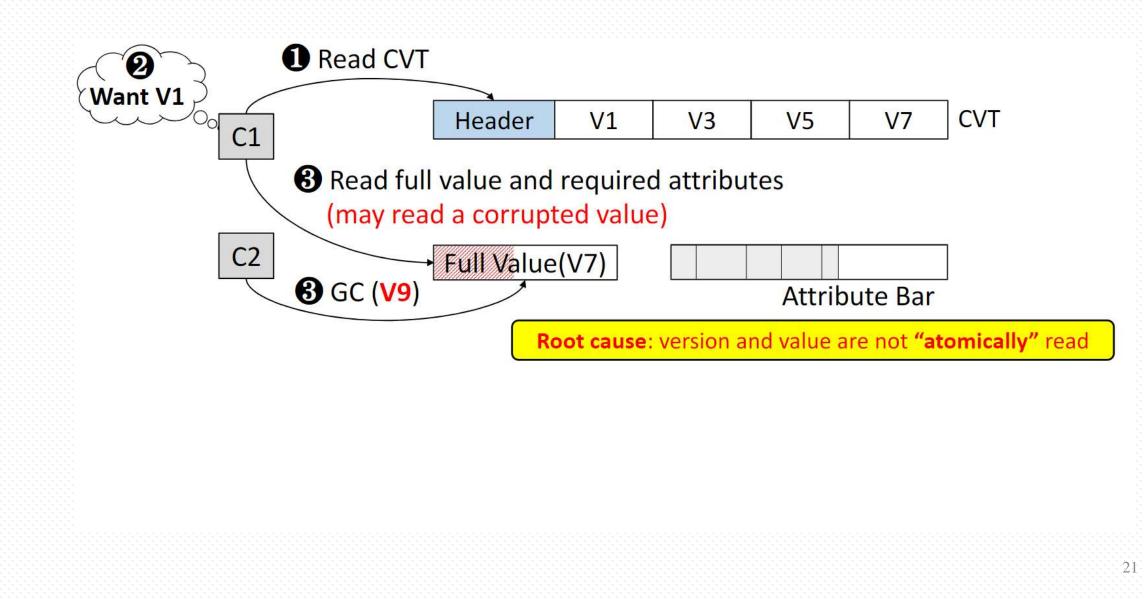
Preer	nptive s	election		
Header	71	V3	V5	V7
		Û		
Header	V9	V3	V5	V7

Overwrite the oldest version Simple, no tracking Low abort rate with fast RDMA USTC, CHINA

nci

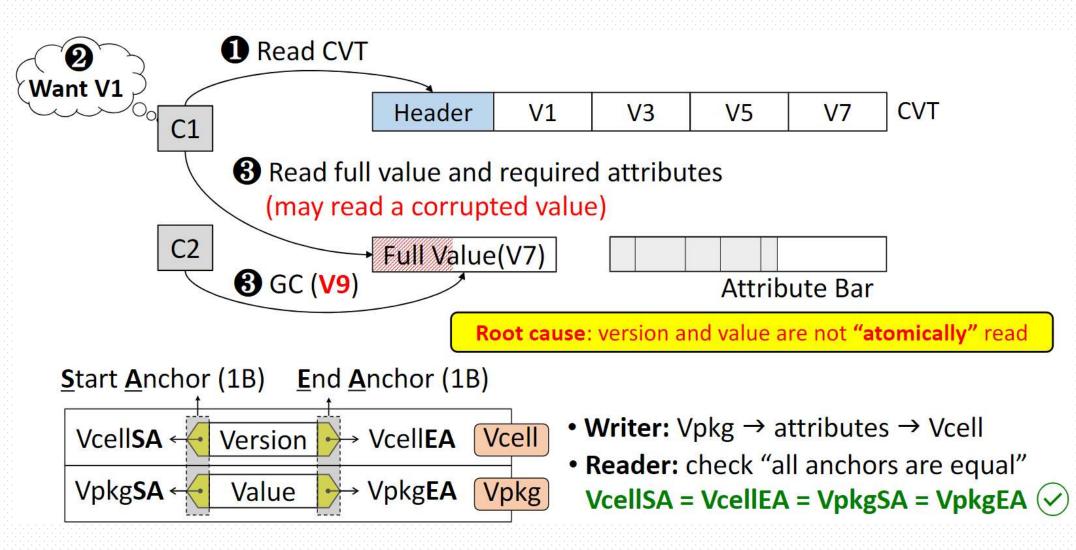
Anchor-Assisted Read





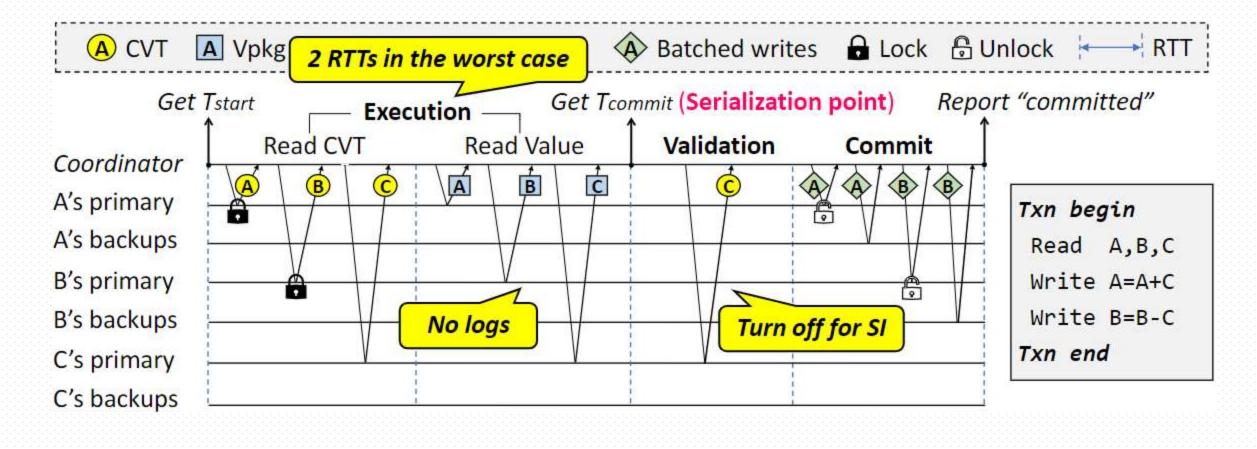
Anchor-Assisted Read





One-Sided RDMA-Based MVCC





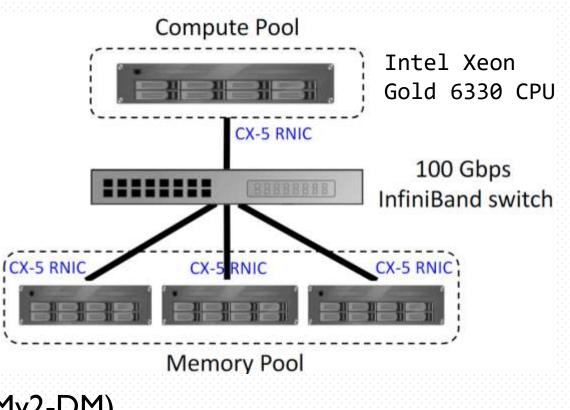
Evaluation

Workloads

- KV Store
 - 8B key + 40B value
 - Skewed (skewness tunable)
- TATP
 - RO/RW: 80%/20%, max 48B
- SmallBank
 - RO/RW: 15%/85%, 16B
- TPCC
 - RO/RW: 8%/92%, max 672B

Comparisons

- FaRMv2@SIGMOD'19 (referred as FaRMv2-DM)
- FORD@FAST'22

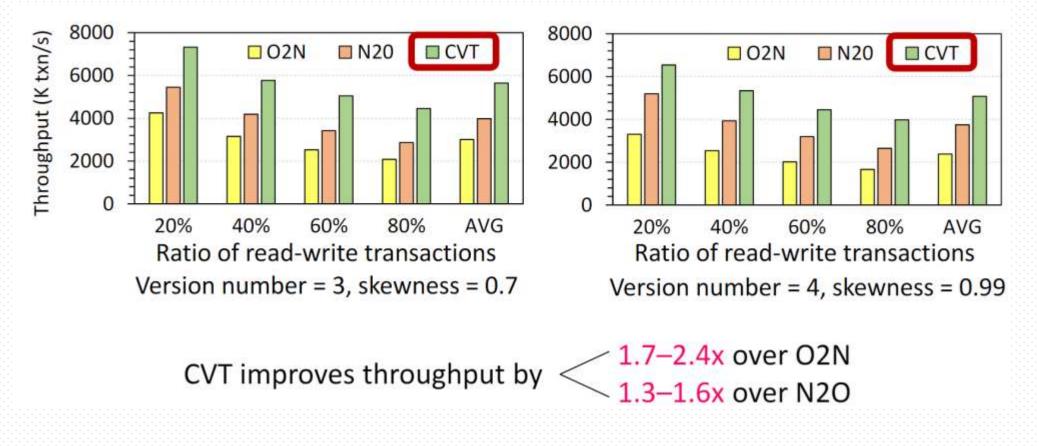




25

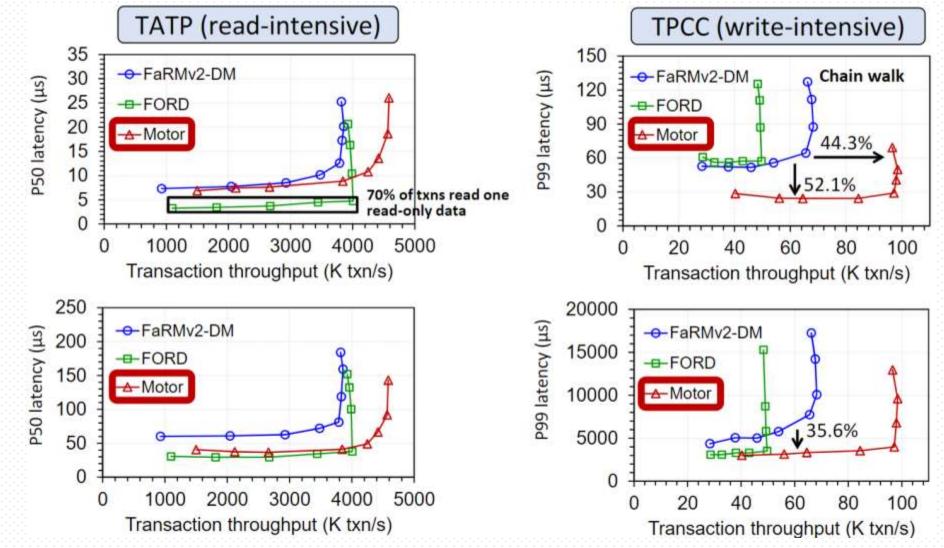
Performance of Version Structures

KV Store



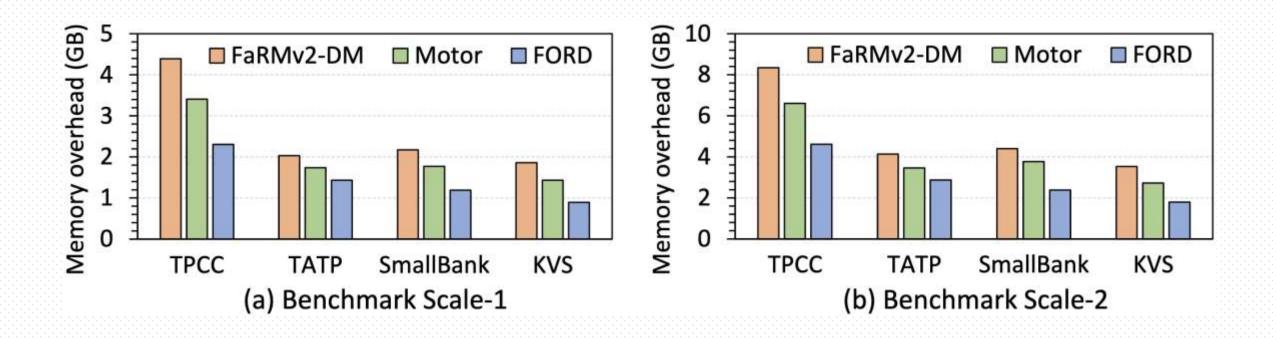
End-to-end Performance





Memory Overhead





Conclusion



Existing multi-version distributed transactions do not fit DM

- Inefficient linked version chain
- Incompatible transaction protocol

Motor: a holistic multi-version design for DM

- Consecutive version tuple structure (memory pool)
- One-sided RDMA MVCC based on CVT (compute pool)

Benefits



https://github.com/minghust/motor



Motor: Enabling Multi-Versioning for Distributed Transactions on

Disaggregated Memory

Ming Zhang, Yu Hua, Zhijun Yang Huazhong University of Science and Technology, China

Thanks you for your attention

